

Chapter 5 Digital Modulation & Detection

5.1 Signal space analysis

Every T seconds, the system sends $k = \log_2 M$ bits.

The rate is $R = k/T$. Each bit sequence of length k

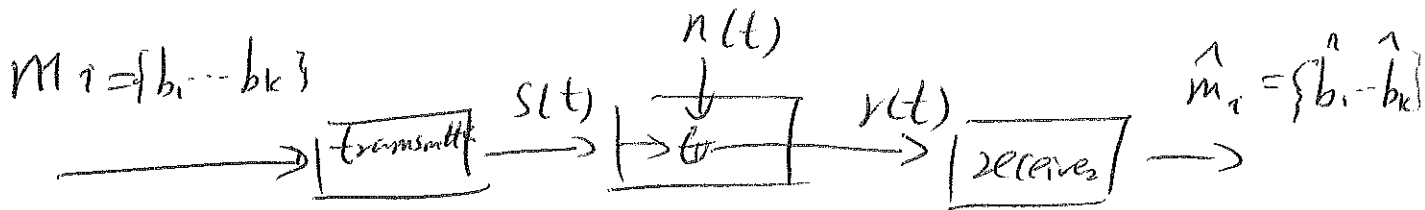
comprise a message $m_i = \{b_1, \dots, b_k\} \in M$, where $M =$

$\{m_1, \dots, m_M\}$ is the set of all such message. ~~the~~ let

The probability P_i ~~of being~~ ^{of being} the selected sequences has $\sum_{i=1}^M P_i = 1$

each message $m_i \in M$ is mapped to a unique analog signal $s_i(t) \in S = \{s_1(t), \dots, s_M(t)\}$, with energy

$$E_{s_i} = \int_0^T s_i^2(t) dt, \quad \forall i$$



where $\hat{m}$ is the best estimate ^{AWGN} ~~and~~ from $r(t) = s(t) + n(t)$

The goal is to minimize the error probability

$$P_e = \sum_{i=1}^M P(\hat{m} \neq m_i | m_i \text{ sent}) P(m_i \text{ sent})$$

5.1.2 Geometric Representation of Signals

orthogonal basis functions $\{\phi_1(t) \dots \phi_N(t)\}_{N \leq M}$

$$s_i(t) = \sum_{j=1}^N s_{ij} \phi_j(t), \quad 0 \leq t < T$$

$$s_{ij} = \int_0^T s_i(t) \phi_j(t) dt$$

s_{ij} is projection of $s_i(t)$ onto $\phi_j(t)$

$$\int_0^T \phi_i(t) \phi_j(t) dt = \begin{cases} 1 & i=j \\ 0 & i \neq j \end{cases}$$

For example $\phi_1(t) = \sqrt{\frac{2}{T}} \cos(2\pi f_c t)$

$$\phi_2(t) = \sqrt{\frac{2}{T}} \sin(2\pi f_c t)$$

$$s_i(t) = s_{i1} \sqrt{\frac{2}{T}} \cos(2\pi f_c t) + s_{i2} \sqrt{\frac{2}{T}} \sin(2\pi f_c t)$$

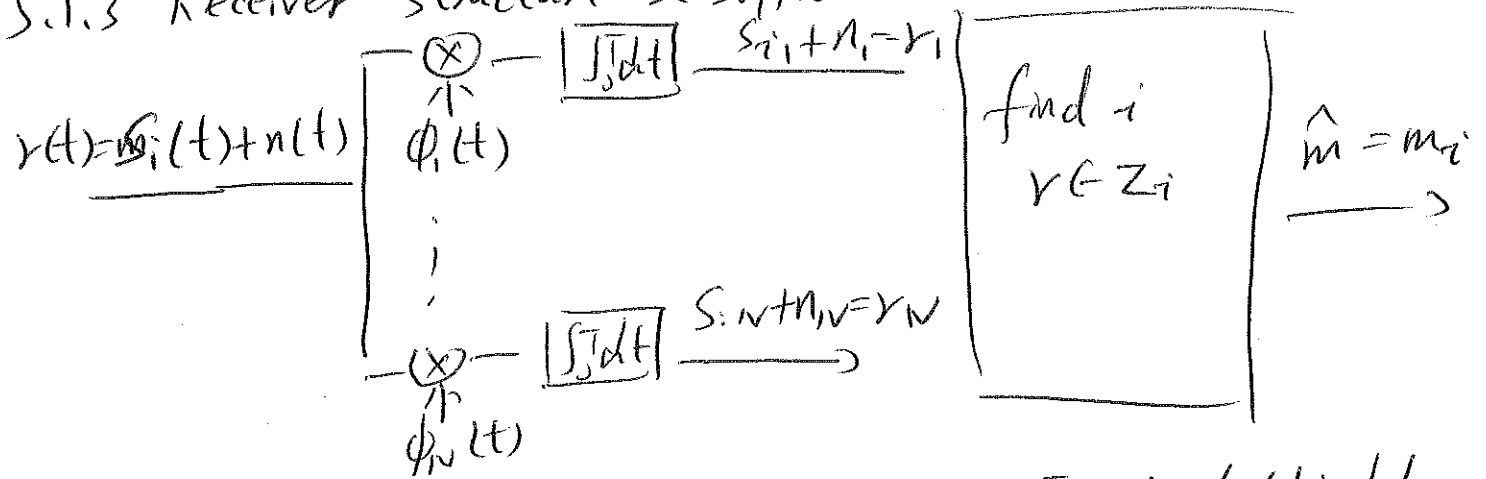
example 5.1 BPSK and QPSK, MQAM

- ^{length}~~power~~ of a vector $\|s_i\| \stackrel{\Delta}{=} \sqrt{\sum_{j=1}^N s_{ij}^2}$

- distance $\|s_i - s_k\| = \sqrt{\sum_{j=1}^N (s_{ij} - s_{kj})^2}$

- ~~int~~ inner product $\langle s_i(t), s_k(t) \rangle = \int_0^T s_i(t) s_k(t) dt$

5.1.3 Receiver structure & sufficient statistics



$$s_{ij} = \int_0^T s_i(t) \phi_j(t) dt, \quad n_j = \int_0^T n(t) \phi_j(t) dt$$

$$r(t) = \sum_{j=1}^N s_j \phi_j(t) + n_r(t)$$

where $s_j = s_{ij} + n_j$ and $n_r(t) = n(t) = \sum_{j=1}^N n_j \phi_j(t)$

We want to see the mean and variance

$$\mu_{r_i | s_i} = E[r_i | s_i] = E[s_{ij} + n_j | s_{ij}] = s_{ij}$$

$$\sigma_{r_i | s_i}^2 = E[r_i - \mu_{r_i | s_i}]^2 = E[n_j^2 | s_{ij}] = E[n_j^2]$$

$$\text{cov}[r_j, r_k | s_i] = \begin{cases} N_0/2 & j=k \\ 0 & j \neq k \end{cases}$$

This is because of AWGN.

we can write probability

we are interested in $P(r | s_i) = \prod_{j=1}^N P(r_j | m_i) = \frac{1}{(\pi N_0)^{N/2}} \exp \left[-\frac{1}{N_0} \sum_{j=1}^N (r_j - s_{ij})^2 \right]$

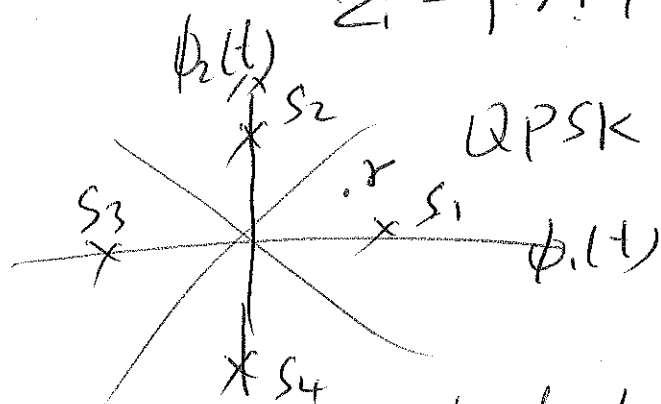
$$P(s_i | r) = \frac{P(r | s_i) P(s_i)}{P(r)} \quad \text{Bayesian rule.}$$

$P(r) \leftarrow$ sufficient statistics

5.1.4 Decision Region and maximum likelihood decision

Decision region $\{Z_1, \dots, Z_M\}$

$$Z_i = \{r: P(s_i | r) > P(s_j | r), \forall j \neq i\}$$



Maximum likelihood decision

$$\arg \max_{s_i} P(s_i | r) = \arg \max_{s_i} \frac{P(r | s_i) P(s_i)}{P(r)} = \arg \max_{s_i} P(r | s_i)$$

if $P(s_i) = P(s_j) \forall i, j$

$$\arg \max_{s_i} \underbrace{P(s_i)}_{=1/M} P(r | s_i), \forall i$$

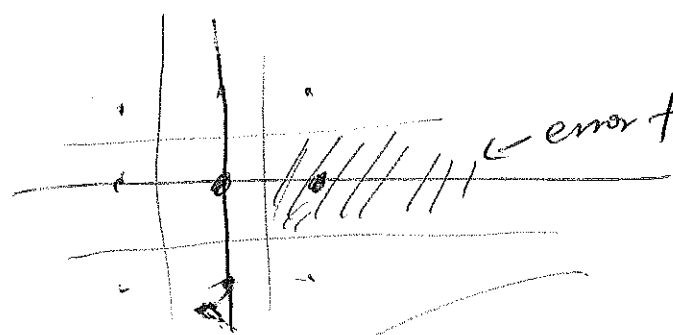
For Gaussian case

$$L(s_i) = -\frac{1}{N_0} \|r - s_i\|^2$$

$$Z_i = \{r: \|r - s_i\| < \|r - s_j\|, \forall j \neq i\}$$

example 5.2

5.1.5 error probability & Union bound



Symbol error rate.

$$P_e(m_i \text{ sent}) = P\left(\bigcup_{\substack{k=1 \\ k \neq i}}^M A_{ik}\right) \leq \sum_{\substack{k=1 \\ k \neq i}}^M P(A_{ik})$$

$$P(A_{ik}) = P(\|s_k - r\| < \|s_i - r\| \mid s_i \text{ sent})$$

$$= P(\|s_k - (s_i + n)\| < \|s_i - (s_i + n)\|)$$

$$= P(\|n + s_i - s_k\| < \|n\|)$$

$$= P(n > \frac{d_{ik}}{2}) = \int_{d_{ik}/2}^{\infty} \frac{1}{\sqrt{2\pi}N_0} \exp\left(-\frac{v^2}{N_0}\right) dv$$

BPSK

$$= Q\left(\frac{d_{ik}}{\sqrt{2N_0}}\right)$$

$$\rightarrow P_e \leq (M-1) Q\left(\frac{d_{\min}}{\sqrt{2N_0}}\right)$$

example 5.3

Gray coding

$$1^2b \leq \frac{P_e}{\log_2 M}$$